

ON THE CROSSING NUMBER OF THE JOIN OF FIVE VERTEX GRAPH G WITH THE DISCRETE GRAPH D_n

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ABSTRACT

In this paper, we show the values of crossing numbers for join products of graph G on five vertices with the discrete graph D_n and the path P_n on n vertices. The proof is done with the help of software. The software generates all cyclic permutations for a given number n . For cyclic permutations, $P_1 - P_m$ will create a graph in which to calculate the distances between all vertices of the graph. These distances are used in proof of crossing numbers of presented graphs.

Keywords: crossing number, cyclic permutations, drawing, graph, join

1. INTRODUCTION

Let G be a simple graph with the vertex set V and the edge set E . A *drawing* of the graph G is a representation of G in the plane such that its vertices are represented by distinct points and its edges by simple continuous arcs connecting the corresponding point pairs. In such a drawing, the intersection of the interiors of the arcs is called a *crossing*. We assume that in a drawing no edge passes through any vertex other than its end-points, no two edges touch each other (i.e., if two edges have a common interior point, then they cross properly at this point), and no three edges cross at the same point. It is easy to see that a drawing with minimum number of crossings (an optimal drawing) is always a *good* drawing, meaning that no edge crosses itself, no two edges cross more than once, and no two edges incident with the same vertex cross each other.

The *crossing number* $cr(G)$ of a simple graph G with the vertex set $V(G)$ and the edge set $E(G)$ is defined as the minimum possible number of edge crossings in a drawing of G in the plane.

Let $G_1 = (V(G_1), E(G_1))$ and $G_2 = (V(G_2), E(G_2))$ be simple graphs. The *join product* of two graphs G_1 and G_2 , denoted by $G_1 + G_2$, is obtained from the vertex-disjoint copies of G_1 and G_2 by adding all edges between $V(G_1)$ and $V(G_2)$. For $|V(G_1)| = m$, and $|V(G_2)| = n$, the edge set of $G_1 + G_2$ is the union of disjoint edge sets of the graphs G_1 , G_2 , and the complete bipartite graph $K_{m,n}$.

Let $D(D(G))$ be a good drawing of the graph G . We denote the number of crossings in D by $cr_D(G)$. Let G_i and G_j be edge-disjoint subgraphs of G . We denote the number of crossings between edges of G_i and edges of G_j by $cr_D(G_i, G_j)$, and the number of crossings among edges of G_i in D by $cr_D(G_i)$. It is easy to see that for any three mutually edge-disjoint subgraphs G_i , G_j , and G_k of G , the following equations hold:

$$cr_D(G_i \cup G_j) = cr_D(G_i) + cr_D(G_j) + cr_D(G_i, G_j),$$

$$cr_D(G_i \cup G_j, G_k) = cr_D(G_i, G_k) + cr_D(G_j, G_k).$$

In the paper, some proofs are based on the Kleitman's result on crossing numbers of complete bipartite graphs. More

precisely, he proved that

$$cr(K_{m,n}) = \left\lfloor \frac{m}{2} \right\rfloor \left\lfloor \frac{m-1}{2} \right\rfloor \left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor, \quad \text{if } m \leq 6.$$

2. SUBJECT

2.1. Software description

We will describe in this subchapter the software which we use when proving the Theorem 4.1 in this article and also in a similar proofs of theorems such like this. The input for the algorithm is the number n , which represents an n -element set $\{1, 2, 3, \dots, n\}$. The algorithm selects all cyclic permutations from the set of all permutations of the n -element set $\{1, 2, 3, \dots, n\}$. The software marks these permutations with symbols P_1, \dots, P_m , where $m = (n-1)!$. Said software gives outputs of distance between each pair of vertices of given graph.

A graph is created with a set of vertices $V = \{P_1, P_2, \dots, P_m\}$ and set of edges E , where the two vertices are joined by the edge if the vertices correspond to the permutations P_i and P_j , which are formed by the exchange of exactly two elements of the n -tuple (i.e. an ordered set with n elements). This graph is represented by a square symmetrical adjacency matrix. The distance between each pair of vertices are calculated using the properties of the cyclic-order graph CO_5 defined in [5].

The software uses the following graph theory: Let us denote $B^{(1)}$ the matrix is gotten from adjacency matrix B by adding ones to the main diagonal. Let us consider the matrix $B^{(2)} = \{b_{ij}^{(2)}\}_{i,j=1}^m$ such that $B^{(2)} = B^{(1)} \cdot B^{(1)}$. From the matrix multiplication it is obvious that $b_{ij}^{(2)} = \sum_{k=1}^m b_{ik}^{(1)} \cdot b_{kj}^{(1)}$, but in this matrix we will use the Boolean addition and multiplication ($1 \cdot 1 = 1$, $0 \cdot 1 = 1 \cdot 0 = 0 \cdot 0 = 0$, $1 + 0 = 0 + 1 = 1 + 1 = 1$, and $0 + 0 = 0$). Generally we can consider matrix $B^{(m)} = B^{(m-1)} \cdot B^{(1)}$.

Theorem 2.1. *Let the adjacency matrix B of the connected graph $G = (V, H)$, $|V| = n$ is given. Then for arbitrary $k = 1, 2, \dots, m$, the element $b_{ij}^{(k)}$ of the matrix $B^{(k)}$ is equal to one if $d(v_i, v_j) \leq k$.*

Corollary 2.1. *The graph $G = (V, H)$, $|V| = n$, is con-*

nected only when the elements of the matrix $B^{(n-1)}$ are only ones.

Corollary 2.2. For each two different vertices of the graph $G = (V, H)$, $d(v_i, v_j) = \min_{k \in \{1, 2, \dots, m\}} \{k; b_{ij}^{(k)} = 1\}$.

2.2. Important facts

We will show the correct proof of the theorem from the article [1]. We deal with the graph G with the vertex set $V = \{v_1, v_2, v_3, v_4, v_5\}$ which is shown in Fig. 1. There is also the graph G with renamed vertices $V = \{1, 2, 3, 4, 5\}$, which is done counter-clockwise with the beginning in the upper right corner.

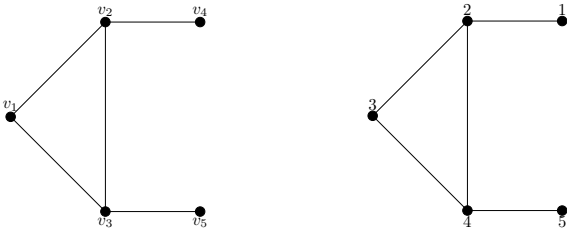


Fig. 1 Five vertex graph G and its numbering of vertices

Using the given software for $n = 5$, we get the names of the cyclic permutations (see Table 1) and the distance between the cyclic permutations (see Table 2-5). The maximum distance between two vertices in the graph is equal to four. With this information we can use arguments which are described in [3,4].

Table 1 Names of Cyclic Permutations of 5-elements set

Name	Cyclic perm.	Name	Cyclic perm.
$P_1 \rightarrow$	(1 2 3 4 5)	$P_{13} \rightarrow$	(1 2 5 4 3)
$P_2 \rightarrow$	(1 3 2 4 5)	$P_{14} \rightarrow$	(1 5 2 4 3)
$P_3 \rightarrow$	(1 2 4 3 5)	$P_{15} \rightarrow$	(1 2 4 5 3)
$P_4 \rightarrow$	(1 4 2 3 5)	$P_{16} \rightarrow$	(1 4 2 5 3)
$P_5 \rightarrow$	(1 4 3 2 5)	$P_{17} \rightarrow$	(1 4 5 2 3)
$P_6 \rightarrow$	(1 3 4 2 5)	$P_{18} \rightarrow$	(1 5 4 2 3)
$P_7 \rightarrow$	(1 2 3 5 4)	$P_{19} \rightarrow$	(1 5 3 4 2)
$P_8 \rightarrow$	(1 3 2 5 4)	$P_{20} \rightarrow$	(1 3 5 4 2)
$P_9 \rightarrow$	(1 2 5 3 4)	$P_{21} \rightarrow$	(1 5 4 3 2)
$P_{10} \rightarrow$	(1 5 2 3 4)	$P_{22} \rightarrow$	(1 4 5 3 2)
$P_{11} \rightarrow$	(1 5 3 2 4)	$P_{23} \rightarrow$	(1 4 3 5 2)
$P_{12} \rightarrow$	(1 3 5 2 4)	$P_{24} \rightarrow$	(1 3 4 5 2)

3. METHODS

We consider the join of G with the discrete graph on n vertices D_n . The graph $G + D_n$ consists of one copy of the graph G and of n vertices t_1, t_2, \dots, t_n , where any vertex $t_i, i = 1, 2, \dots, n$, is adjacent to every vertex of G . Let $T^i, 1 \leq i \leq n$, denote the subgraph induced by the five edges incident with the vertex t_i . Then

$$G + D_n = G \cup K_{5,n} = G \cup \left(\bigcup_{i=1}^n T^i \right).$$

The graph $G + D_1$ is planar, thus $cr(G + D_1) = 0$. One can easily verify that $cr(G + D_2) \leq 1$. The graph $G + D_2$ contains a subdivision of $K_{3,3}$ as a subgraph, and therefore $cr(G + D_2) \geq 1$. So, $cr(G + D_2) = 1$.

Let D be a good drawing of the graph $G + D_n$. The rotation $rot_D(t_i)$ of vertex t_i in the drawing D is the cyclic permutation that records the (cyclic) counter-clockwise order in which the edges leave t_i , see [4]. We use the notation (12543) if the counter-clockwise order of the edges incident with the vertex t_i is t_i1, t_i2, t_i5, t_i4 , and t_i3 (see A_1 in Fig. 2), where $V = \{1, 2, 3, 4, 5\}$ are noted vertices of the graph G . We emphasize that a rotation is a cyclic permutation. For $i, j \in \{1, 2, \dots, n\}, i \neq j$, every subgraph $T^i \cup T^j$ of the graph $G + D_n$ is isomorphic with the graph $K_{5,2}$. In the paper, we will deal with the minimum necessary number of crossings between the edges of T^i and the edges of T^j in a subgraph $T^i \cup T^j$ induced by the drawing D of the graph $G + D_n$ depending on the rotations $rot_D(t_i)$ and $rot_D(t_j)$.

D. R. Woodall [5] proved that, in any good drawing D of the graph $K_{5,2}, cr_D(T^i, T^j) \geq 4$ if $rot_D(t_i) = rot_D(t_j)$. Moreover, if $Q(rot_D(t_i), rot_D(t_j))$ denotes the minimum number of interchanges of adjacent elements of $rot_D(t_i)$ required to produce the inverse cyclic permutation of $rot_D(t_j)$, then $Q(rot_D(t_i), rot_D(t_j)) \leq cr_D(T^i, T^j)$.

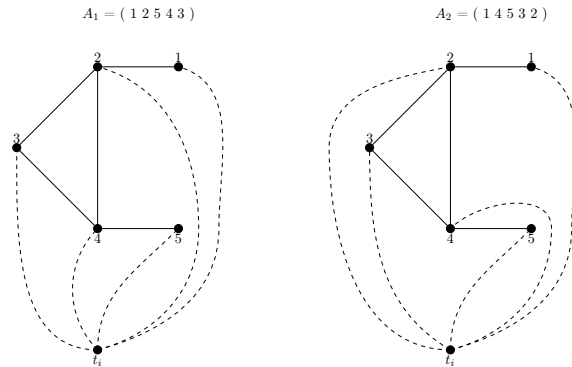


Fig. 2 Graph G and its configurations - type A

We will separate the subgraphs $T^i, i = 1, \dots, n$, of $G + D_n$ into three subsets depending on how many the considered T^i crosses the edges of G in D . For $i = 1, 2, \dots, n$, we denote by $R_D = \{T^i : cr_D(G, T^i) = 0\}$ and $S_D = \{T^i : cr_D(G, T^i) = 1\}$. Every other subgraph T^i crosses G at least twice in D . Moreover, let F^i denote the subgraph $G \cup T^i$ for $T^i \in R_D$, where $i \in \{1, \dots, n\}$. Thus, any F^i is exactly represented by $rot_D(t_i)$.

Table 2 Distance 1 Between Cyclic Permutations

From	To
P_1	$P_2, P_3, P_7, P_{10}, P_{24}$
P_2	$P_1, P_6, P_8, P_{11}, P_{15}$
P_3	$P_1, P_4, P_{14}, P_{15}, P_{23}$
P_4	$P_3, P_5, P_7, P_{16}, P_{18}$
P_5	$P_4, P_6, P_8, P_{21}, P_{23}$
P_6	$P_2, P_5, P_{16}, P_{19}, P_{24}$
P_7	$P_1, P_4, P_8, P_9, P_{20}$
P_8	$P_2, P_5, P_7, P_{12}, P_{13}$
P_9	$P_7, P_{10}, P_{13}, P_{16}, P_{19}$
P_{10}	$P_1, P_9, P_{11}, P_{14}, P_{17}$
P_{11}	$P_2, P_{10}, P_{12}, P_{19}, P_{22}$
P_{12}	$P_8, P_{11}, P_{14}, P_{20}, P_{23}$
P_{13}	$P_8, P_9, P_{14}, P_{15}, P_{21}$
P_{14}	$P_3, P_{10}, P_{12}, P_{13}, P_{18}$
P_{15}	$P_2, P_3, P_{13}, P_{16}, P_{22}$
P_{16}	$P_4, P_6, P_9, P_{15}, P_{17}$
P_{17}	$P_{10}, P_{16}, P_{18}, P_{22}, P_{24}$
P_{18}	$P_4, P_{14}, P_{17}, P_{20}, P_{21}$
P_{19}	$P_6, P_9, P_{11}, P_{20}, P_{21}$
P_{20}	$P_7, P_{12}, P_{18}, P_{19}, P_{24}$
P_{21}	$P_5, P_{13}, P_{18}, P_{19}, P_{22}$
P_{22}	$P_{11}, P_{15}, P_{17}, P_{21}, P_{23}$
P_{23}	$P_3, P_5, P_{12}, P_{22}, P_{24}$
P_{24}	$P_1, P_6, P_{17}, P_{20}, P_{23}$

Table 4 Distance 3 Between Cyclic Permutations

From	To
P_1	$P_5, P_{12}, P_{13}, P_{16}, P_{18}, P_{19}, P_{22}$
P_2	$P_4, P_9, P_{14}, P_{17}, P_{20}, P_{21}, P_{23}$
P_3	$P_6, P_8, P_9, P_{11}, P_{17}, P_{20}, P_{21}$
P_4	$P_2, P_{10}, P_{12}, P_{13}, P_{19}, P_{22}, P_{24}$
P_5	$P_1, P_9, P_{11}, P_{14}, P_{15}, P_{17}, P_{20}$
P_6	$P_3, P_7, P_{10}, P_{12}, P_{13}, P_{18}, P_{22}$
P_7	$P_6, P_{11}, P_{14}, P_{15}, P_{17}, P_{21}, P_{23}$
P_8	$P_3, P_{10}, P_{16}, P_{18}, P_{19}, P_{22}, P_{24}$
P_9	$P_2, P_3, P_5, P_{12}, P_{18}, P_{22}, P_{24}$
P_{10}	$P_4, P_6, P_8, P_{15}, P_{20}, P_{21}, P_{23}$
P_{11}	$P_3, P_5, P_7, P_{13}, P_{16}, P_{18}, P_{24}$
P_{12}	$P_1, P_4, P_6, P_9, P_{15}, P_{17}, P_{21}$
P_{13}	$P_1, P_4, P_6, P_{11}, P_{17}, P_{20}, P_{23}$
P_{14}	$P_2, P_5, P_7, P_{16}, P_{19}, P_{22}, P_{24}$
P_{15}	$P_5, P_7, P_{10}, P_{12}, P_{18}, P_{19}, P_{24}$
P_{16}	$P_1, P_8, P_{11}, P_{14}, P_{20}, P_{21}, P_{23}$
P_{17}	$P_2, P_3, P_5, P_7, P_{12}, P_{13}, P_{19}$
P_{18}	$P_1, P_6, P_8, P_9, P_{11}, P_{15}, P_{23}$
P_{19}	$P_1, P_4, P_8, P_{14}, P_{15}, P_{17}, P_{23}$
P_{20}	$P_2, P_3, P_5, P_{10}, P_{13}, P_{16}, P_{22}$
P_{21}	$P_2, P_3, P_7, P_{10}, P_{12}, P_{16}, P_{24}$
P_{22}	$P_1, P_4, P_6, P_8, P_9, P_{14}, P_{20}$
P_{23}	$P_2, P_7, P_{10}, P_{13}, P_{16}, P_{18}, P_{19}$
P_{24}	$P_4, P_8, P_9, P_{11}, P_{14}, P_{15}, P_{21}$

Table 3 Distance 2 Between Cyclic Permutations

From	To
P_1	$P_4, P_6, P_8, P_9, P_{11}, P_{14}, P_{15}, P_{17}, P_{20}, P_{23}$
P_2	$P_3, P_5, P_7, P_{10}, P_{12}, P_{13}, P_{16}, P_{19}, P_{22}, P_{24}$
P_3	$P_2, P_5, P_7, P_{10}, P_{12}, P_{13}, P_{16}, P_{18}, P_{22}, P_{24}$
P_4	$P_1, P_6, P_8, P_9, P_{14}, P_{15}, P_{17}, P_{20}, P_{21}, P_{23}$
P_5	$P_2, P_3, P_7, P_{12}, P_{13}, P_{16}, P_{18}, P_{19}, P_{22}, P_{24}$
P_6	$P_1, P_4, P_8, P_9, P_{11}, P_{15}, P_{17}, P_{20}, P_{21}, P_{23}$
P_7	$P_2, P_3, P_5, P_{10}, P_{12}, P_{13}, P_{16}, P_{18}, P_{19}, P_{24}$
P_8	$P_1, P_4, P_6, P_9, P_{11}, P_{14}, P_{15}, P_{20}, P_{21}, P_{23}$
P_9	$P_1, P_4, P_6, P_8, P_{11}, P_{14}, P_{15}, P_{17}, P_{20}, P_{21}$
P_{10}	$P_2, P_3, P_7, P_{12}, P_{13}, P_{16}, P_{18}, P_{19}, P_{22}, P_{24}$
P_{11}	$P_1, P_6, P_8, P_9, P_{14}, P_{15}, P_{17}, P_{20}, P_{21}, P_{23}$
P_{12}	$P_2, P_3, P_5, P_7, P_{10}, P_{13}, P_{18}, P_{19}, P_{22}, P_{24}$
P_{13}	$P_2, P_3, P_5, P_7, P_{10}, P_{12}, P_{16}, P_{18}, P_{19}, P_{22}$
P_{14}	$P_1, P_4, P_8, P_9, P_{11}, P_{15}, P_{17}, P_{20}, P_{21}, P_{23}$
P_{15}	$P_1, P_4, P_6, P_8, P_9, P_{11}, P_{14}, P_{17}, P_{21}, P_{23}$
P_{16}	$P_2, P_3, P_5, P_7, P_{10}, P_{13}, P_{18}, P_{19}, P_{22}, P_{24}$
P_{17}	$P_1, P_4, P_6, P_9, P_{11}, P_{14}, P_{15}, P_{20}, P_{21}, P_{23}$
P_{18}	$P_3, P_5, P_7, P_{10}, P_{12}, P_{13}, P_{16}, P_{19}, P_{22}, P_{24}$
P_{19}	$P_2, P_5, P_7, P_{10}, P_{12}, P_{13}, P_{16}, P_{18}, P_{22}, P_{24}$
P_{20}	$P_1, P_4, P_6, P_8, P_9, P_{11}, P_{14}, P_{17}, P_{21}, P_{23}$
P_{21}	$P_4, P_6, P_8, P_9, P_{11}, P_{14}, P_{15}, P_{17}, P_{20}, P_{23}$
P_{22}	$P_2, P_3, P_5, P_{10}, P_{12}, P_{13}, P_{16}, P_{18}, P_{19}, P_{24}$
P_{23}	$P_1, P_4, P_6, P_8, P_{11}, P_{14}, P_{15}, P_{17}, P_{20}, P_{21}$
P_{24}	$P_2, P_3, P_5, P_7, P_{10}, P_{12}, P_{16}, P_{18}, P_{19}, P_{22}$

Table 5 Distance 4 Between Cyclic Permutations

From	To
P_1	P_{21}
P_2	P_{18}
P_3	P_{19}
P_4	P_{11}
P_5	P_{10}
P_6	P_{14}
P_7	P_{22}
P_8	P_{17}
P_9	P_{23}
P_{10}	P_5
P_{11}	P_4
P_{12}	P_{16}
P_{13}	P_{24}
P_{14}	P_6
P_{15}	P_{20}
P_{16}	P_{12}
P_{17}	P_8
P_{18}	P_2
P_{19}	P_3
P_{21}	P_1
P_{22}	P_7
P_{23}	P_9
P_{24}	P_{13}

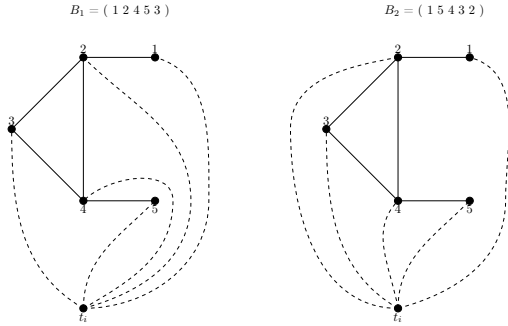


Fig. 3 Graph G and its configurations - type B

There is only one drawing of G without crossings shown in Fig. 1. Assume a good drawing D of the graph $G + D_n$ in which the edges of G does not cross each other. We will count the number of necessary crossings between two subgraphs T^i and T^j with $cr_D(G, T^i \cup T^j) = 0$. In this case, without loss of generality, we can choose the vertex notations of the graph in such a way as shown in Fig. 1. It is easy to see that, in D , there are only four different possible configurations of F^i summarized in Table 6, see Fig. 2 and 3. We denote by \mathcal{M}_D the set of all configurations that exist in the drawing D belonging to the set \mathcal{M} , where $\mathcal{M} = \{A_1, A_2, B_1, B_2\}$.

Table 6 Configurations of graph $G \cup T^i$ with vertices denoted of G as in Fig. 1

$A_1 : (12543)$	$A_2 : (14532)$
$B_1 : (12453)$	$B_2 : (15432)$

Let X, Y be configurations from \mathcal{M}_D . We shortly denote by $cr_D(X, Y)$ the number of crossings in D between T^i and T^j for different $T^i, T^j \in R_D$ such that F^i, F^j have configurations X, Y , respectively. Finally, let $cr(X, Y) = \min\{cr_D(X, Y)\}$ over all good drawings D of the graph $G + D_n$ with $X, Y \in \mathcal{M}_D$. The configuration A_1 is represented by the cyclic permutation $P_{13} = (12543)$ and the configuration A_2 is represented by the cyclic permutation $P_{22} = (14532)$. As $P_7 = (12354)$ is the inverse cyclic permutation to the permutation P_{22} , then $cr(A_1, A_2) \geq 2$ by Table 3. The similar idea is used for the another cases. Thus, all lower-bounds of numbers of crossings of configurations from \mathcal{M} are summarized in Table 7.

Table 7 Lower-bounds of numbers of crossings of two configurations from \mathcal{M}

	A_1	A_2	B_1	B_2
A_1	4	2	3	3
A_2	2	4	3	3
B_1	3	3	4	2
B_2	3	3	2	4

4. RESULTS

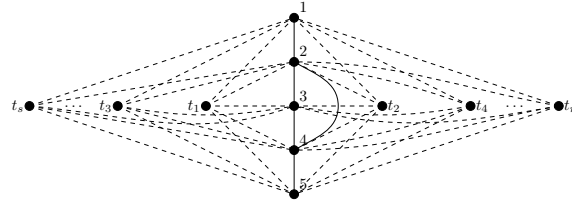


Fig. 4 Good drawing of $G + D_n$

Theorem 4.1. Let G be the graph in Fig. 1 and D_n is discrete graph with n vertices, then

$$cr(G + D_n) = 4 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + \lfloor \frac{n}{2} \rfloor, \text{ for } n \geq 1.$$

Proof: The theorem is true for $n = 1$ and $n = 2$. In Fig. 4 there is a drawing of $G + D_n$ with $4 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + \lfloor \frac{n}{2} \rfloor$ crossings. Thus, $cr(G + D_n) \leq 4 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + \lfloor \frac{n}{2} \rfloor$. We prove the reverse inequality by induction on n . For $n \geq 3$, let D be a good drawing of $G + D_n$ with less than $4 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + \lfloor \frac{n}{2} \rfloor$ crossings. Suppose now that, for $n \geq 3$

$$cr(G + D_{n-2}) \geq 4 \lfloor \frac{n-2}{2} \rfloor \lfloor \frac{n-3}{2} \rfloor + \lfloor \frac{n-2}{2} \rfloor$$

and consider such a drawing D of $G + D_n$ that

$$cr_D(G + D_n) < 4 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + \lfloor \frac{n}{2} \rfloor. \tag{1}$$

The drawing D has the following property:

$$cr_D(T^i, T^j) \neq 0 \quad \text{for all } i, j = 1, 2, \dots, n, i \neq j. \tag{2}$$

To prove it assume that there are two different subgraphs T^i and T^j such that $cr_D(T^i, T^j) = 0$ and let for every integer $s, s < n$, any good drawing of graph $G + D_s$ has at least $4 \lfloor \frac{s}{2} \rfloor \lfloor \frac{s-1}{2} \rfloor + \lfloor \frac{s}{2} \rfloor$ crossings. Without loss of generality let $cr_D(T^{n-1}, T^n) = 0$, one can easy to verify that $cr_D(G, T^{n-1} \cup T^n) \geq 1$. By $cr(K_{5,3}) = 4$ we give $cr_D(T^k, T^{n-1} \cup T^n) \geq 4$ for $k = 1, 2, \dots, n-2$. So, for the number of crossings in D we have

$$\begin{aligned} cr_D(G + D_n) &= cr_D \left(G \cup \bigcup_{i=1}^{n-2} T^i \right) + cr_D(T^{n-1} \cup T^n) + \\ &+ cr_D(G, T^{n-1} \cup T^n) + cr_D \left(\bigcup_{i=1}^{n-2} T^i, T^{n-1} \cup T^n \right) \geq \\ &\geq 4 \lfloor \frac{n-2}{2} \rfloor \lfloor \frac{n-3}{2} \rfloor + \lfloor \frac{n-2}{2} \rfloor + 1 + 4(n-2) = \\ &= 4 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + \lfloor \frac{n}{2} \rfloor. \end{aligned}$$

This contradicts (1), and therefore $cr_D(T^i, T^j) \neq 0$ for all $i, j = 1, 2, \dots, n, i \neq j$. Our assumption on D together with $cr(K_{5,n}) = 4 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor$ implies that

$$cr_D(G) + cr_D(G, K_{5,n}) < \lfloor \frac{n}{2} \rfloor.$$

Thus, we have $r = |R_D| > \lfloor \frac{n}{2} \rfloor, s = |S_D| < \lfloor \frac{n}{2} \rfloor$.

Case 1: $cr_D(G) = 0$ as in Fig. 1.

- a) $\{A_1, A_2\} \subseteq \mathcal{M}_D$ or $\{B_1, B_2\} \subseteq \mathcal{M}_D$.
 Without loss of generality if we fix any two $T^n, T^{n-1} \in R_D$ such that F^n, F^{n-1} have configurations A_1, A_2 , respectively, then $cr_D(G \cup T^n \cup T^{n-1}, T^i) \geq 6$ holds by Table 7 for any $T^i \in R_D$. Using (2) we have

$$\begin{aligned} cr_D(G + D_n) &= cr_D(K_{5,n-2}) + cr_D(G \cup T^n \cup T^{n-1}) + \\ &+ cr_D(K_{5,n-2}, G \cup T^n \cup T^{n-1}) \geq 4 \lfloor \frac{n-2}{2} \rfloor \lfloor \frac{n-3}{2} \rfloor + \\ &+ 6(r-2) + 3s + 4(n-r-s) + 2 = \\ &= 4 \lfloor \frac{n-2}{2} \rfloor \lfloor \frac{n-3}{2} \rfloor + 2r - s + 4n - 10 \geq \\ &\geq 4 \lfloor \frac{n-2}{2} \rfloor \lfloor \frac{n-3}{2} \rfloor + 2 \left(\lfloor \frac{n}{2} \rfloor + 1 \right) + \\ &+ 1 - \lfloor \frac{n}{2} \rfloor + 4n - 10 \geq 4 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + \lfloor \frac{n}{2} \rfloor. \end{aligned}$$

- b) $\{A_1, A_2\} \not\subseteq \mathcal{M}_D$ and $\{B_1, B_2\} \not\subseteq \mathcal{M}_D$.
 Without loss of generality if we fix any $T^n \in R_D$ such that F^n has configuration from \mathcal{M}_D , then $cr_D(G \cup T^n, T^i) \geq 3$ holds for any $T^i \in R_D$. Using the property (2) we have

$$\begin{aligned} cr_D(G + D_n) &= cr_D(K_{5,n-1}) + cr_D(G \cup T^n) + \\ &+ cr_D(K_{5,n-1}, G \cup T^n) \geq 4 \lfloor \frac{n-1}{2} \rfloor \lfloor \frac{n-2}{2} \rfloor + \\ &+ 3(r-1) + 2s + 3(n-r-s) = \\ &= 4 \lfloor \frac{n-1}{2} \rfloor \lfloor \frac{n-2}{2} \rfloor + 3n - s - 3 \geq \\ &\geq 4 \lfloor \frac{n-1}{2} \rfloor \lfloor \frac{n-2}{2} \rfloor + 3n + 1 - \lfloor \frac{n}{2} \rfloor - 3 \geq \\ &\geq 4 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + \lfloor \frac{n}{2} \rfloor. \end{aligned}$$

In the following three cases we will use the same idea as in the Case 1 b). Since $r = |R_D| > \lfloor \frac{n}{2} \rfloor$, the vertices of degree one noted by 1, 5 cannot be separated by 3-cycle of the graph G .

Case 2: $cr_D(G) = 1$ as in Fig. 5.

By a discussion we can easily verify that there is only one type of configuration for F^i represented by $P_5 = (14325)$.

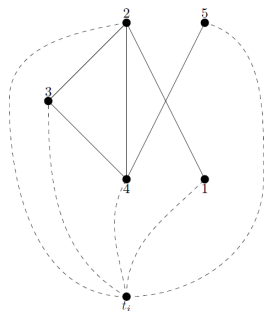


Fig. 5 Good drawing of $G \cup T^i$

Case 3: $cr_D(G) = 1$ as in Fig. 6.

By a discussion we can verify that there are only two types of configurations for F^i represented by the cyclic permutations $P_{23} = (14352)$ and $P_3 = (12435)$. $P_{19} = (15342)$ is the inverse cyclic permutation to the permutation P_3 . Thus, by Table 4 we give lower-bound of number of crossings of these configurations equal to three.

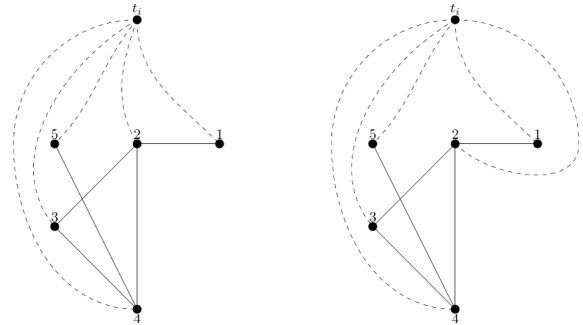


Fig. 6 Good drawing of $G \cup T^i$

Case 4: $cr_D(G) = 3$ as in Fig. 7.

By a discussion we can verify that there is only one type of configuration for F^i represented by the cyclic permutation $P_{12} = (13524)$.

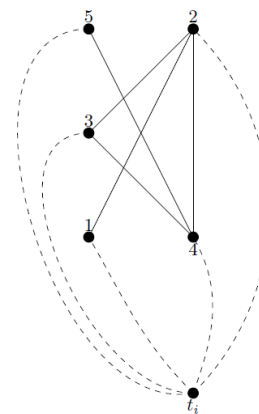


Fig. 7 Good drawing of $G \cup T^i$

Theorem 4.2. Let G be the graph in Fig. 1 and P_n is a path on n vertices, then

$$cr(G + P_n) = 4 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + \lfloor \frac{n}{2} \rfloor \text{ for } n \geq 1.$$

We are able to add the edges without crossings in Fig. 4. So the drawing of the graph $G + P_n$ with $4 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + \lfloor \frac{n}{2} \rfloor$ crossings is obtained.

5. CONCLUSIONS

In this article, we show the proof technique for a crossing number in a given graphs that used the data generated by the software. More significant usage of this software occurs for larger values of n than five. We get 120 cyclical permutations for $n = 6$, 720 cyclical permutations for $n = 7$ which is significantly more than 24 cyclical permutations for $n = 5$. For such values, software is an indispensable tool since, we get considerably more complicated graph of distances between cyclic permutations.

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